Timer Tosks

ACM SIGCOMM August 19-24, 2019 Beijing, China

Towards Time-driven Execution in Programmable Dataplanes

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Problem

Programmable dataplanes provide event-driven execution

- By means of the match-action paradigm
- Event = arrival of a packet OR a specific packet (match specification)

Orchestrating TimerTasks

Challenge

- Time-driven abstraction on top of packet-driven hardware

Any programmed logic executes when there is a packet

- But if there is no packet, **nothing happens**! _
- Not a problem for implementing flexible forwarding pipelines

Increasingly used for advanced applications

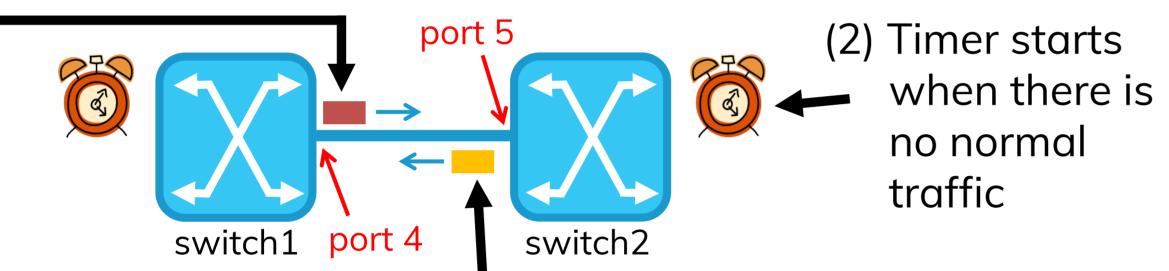
- In-network protocols
- Distributed systems that are co-designed with the network

A purely event-driven execution model is insufficient!

Need for Time-driven Execution

Fast and efficient link failure detection in the dataplane

(1) Packets from normal traffic act as "heartbeats"



(3) On each timeout, explicit ping packets are sent to check connectivity

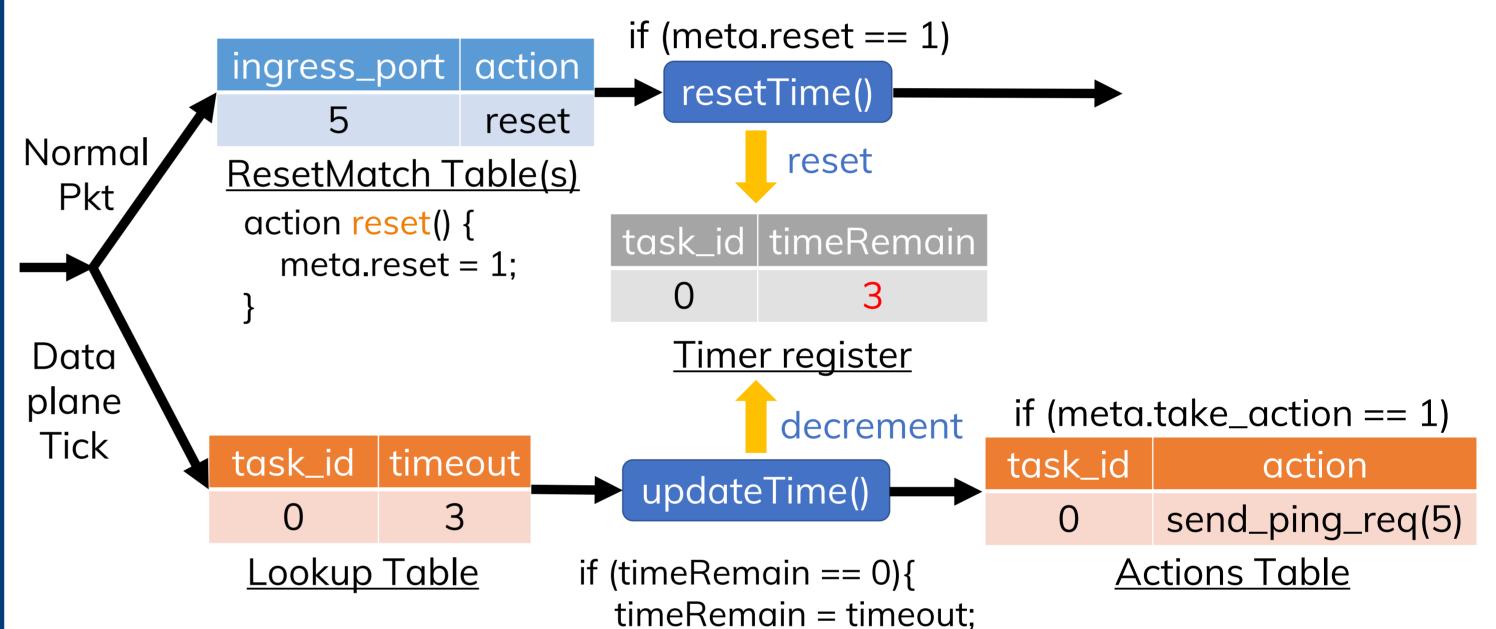
Solution approach: inspired from Linux OS

- Linux is event-driven
- Provides time-driven execution via periodic events (kernel ticks)

"Periodic-Event Framework" for the dataplane

- Dataplane ticks: periodic and regular packets
- On-chip hardware packet generator: source of dataplane ticks —
- Orchestrate TimerTasks: decrement and reset timers, perform actions

<u>TimerTask link_failure_detection for switch2 (port 5)</u>



Not possible to achieve with a packet-driven execution model \rightarrow An action is required in response to the <u>absence</u> of a packet

Avoid aging metrics, refresh them!

- In CONGA [SIGCOMM '14] and DistCache [FAST '19], source ToR blindly ages the metrics when there is no reverse traffic
- Instead, destination ToR can explicitly refresh the metrics in the absence of reverse traffic

Mechanisms for reliable message delivery in the dataplane

- Msg retransmission in dataplane in the <u>absence</u> of an ACK (timeout)

Simple Periodic Tasks

- Periodically measure link utilization every 25µs (Zhang et al. [IMC '17])

The TimerTask Abstraction

What kind of execution semantics do we need?

- If (condition == true) for certain timeout \rightarrow execute an action
- e.g. If no heartbeat packet within 100ms, inform the master
- Periodic tasks can be expressed with an always true condition

meta.take_action = 1;

Evaluation Case Studies

Evaluation Setup for Periodic-Event Framework

- Barefoot Tofino switch and two x86 servers in linear topology
- Dataplane ticks: Tofino's on-chip packet generator

Dataplane tick inter-arrival time

- 99^{th} percentile error in tick inter-arrival time (1µs to 100ms) < 0.1%

1) Fast and efficient link failure detection protocol

- Explicitly pings adjacent switch when no packet arrives for 3μ s
- Detects link failures within 6µs

2) Metric refreshing mechanism

- Explicitly updates piggybacked metrics in lieu of normal traffic —
- Future work: quantify application-level benefits

3) High-resolution network measurements

A TimerTask consists of:

- ResetMatch: same as a P4 "match" specification (just different name) —
- Timeout: a positive constant
- Action: standard P4 action

```
timertask link_failure_detection {
resetmatch = {
   standard_metadata.ingress_port: exact;
timeout = 3; // in microseconds
action = send_ping_request;
```

How does it work?

- The system runtime implements an implicit timer
- If a packet matches the ResetMatch specification: the timer is reset
- If the timer times out: the action is executed & the timer is restarted

Periodic task to read switch counters at 1µs interval

Future Work

Target-specific TimerTask P4 Program Standard P4 Target with TimerTasks Compiler Compiler P4 Program

TimerTask Compiler

- Timer start and stop APIs in the dataplane
- Table placement: TimerTasks interact with rest of the P4 program
- Multiplexing: single dataplane tick packet to update multiple timers
- TimerTask runtime (control plane): add/remove TimerTask "instances"

Native TimerTasks in hardware